

AMPyC

Chapter 7: Stochastic Model Predictive Control II

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Authors: Prof. Melanie Zeilinger
Dr. Lukas Hewing
Dr. Andrea Carron
Dr. Johannes Köhler

ETH Zurich

Recap: MPC for additive disturbances - Stochastic setting

Uncertain constrained system

$$x(k+1) = f(x(k), u(k)) + w(k) \quad \Pr(x(k) \in \mathcal{X}) \geq p, \Pr(u(k) \in \mathcal{U}) \geq p, \quad w(k) \sim Q^w, i.i.d.$$

Design control law $u(k) = \pi(x(k))$ such that the system:

1. Satisfies constraints : $x(k) \in \mathcal{X}$, $u(k) \in \mathcal{U}$ **with given probability p**
2. Is 'stable': Converges to the origin **in a suitable sense**
3. Optimizes (nominal/**expected**) "performance"
4. Maximizes the set $\{x(0) \mid \text{Conditions 1-3 are met}\}$

Recap: Asymptotic Average Performance in Stochastic MPC

Asymptotic average performance bound

$$l_{\text{avg}} = \lim_{\bar{N} \rightarrow \infty} \frac{1}{\bar{N}} \sum_{k=0}^{\bar{N}-1} \mathbb{E}(l(x(k), \pi(x(k)))) \leq C$$

Derive such a performance bound for stochastic MPC in three steps:

1. Lyapunov-like decrease implies asymptotic average performance bound

$$\mathbb{E}(V(x(k+1)) | x(k)) - V(x(k)) \leq -l(x(k), \pi(x(k))) + C \Rightarrow l_{\text{avg}} \leq C$$

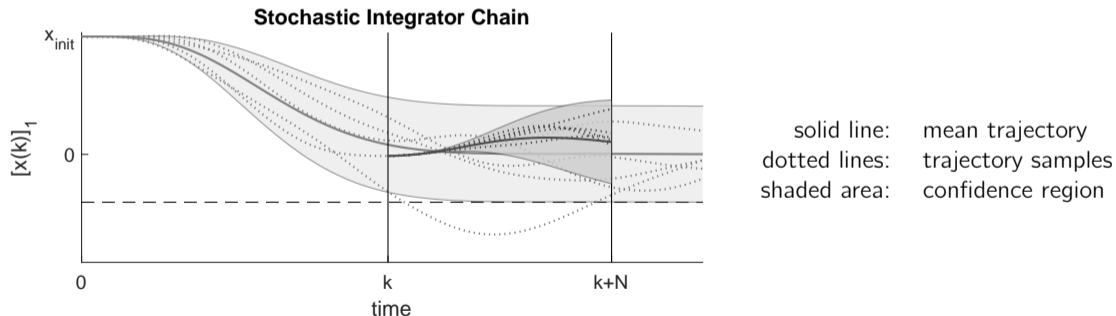
2. Apply this to the cost function decrease in MPC

$$\mathbb{E}(J^*(x(k+1)) | x(k)) - J^*(x(k)) \leq -l(x(k), \pi(x(k))) + C$$

3. Tractable solution for (unconstrained/recursively feasible) linear stochastic MPC

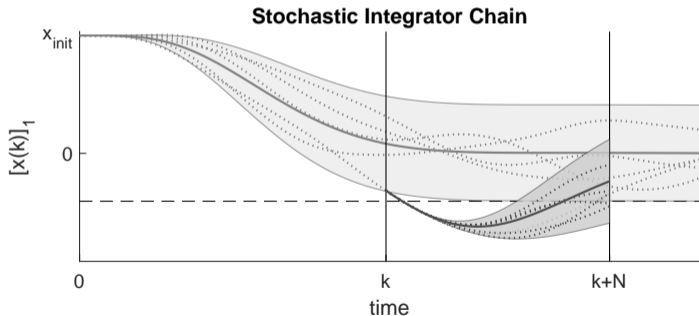
Recap: Feasibility Issues in SMPC

Under general stochastic disturbances, feasibility can usually **not** be guaranteed



Recap: Feasibility Issues in SMPC

Under general stochastic disturbances, feasibility can usually **not** be guaranteed



solid line: mean trajectory
dotted lines: trajectory samples
shaded area: confidence region

Recap: Closed-loop chance constraint satisfaction

Closed-loop chance constraints:

$$(*) \Pr(x(k) \in \mathcal{X} \mid x(0)) \geq p, \forall k \geq 0$$

But MPC formulation successively enforces

$$(**) \Pr(x(k+1) \in \mathcal{X} \mid x(k)) \geq p, \forall k \geq 0$$

It can be easily seen that $(**) \Rightarrow (*)$ since

$$\Pr(x(k+1) \in \mathcal{X} \mid x(0)) = \int \underbrace{\Pr(x(k+1) \in \mathcal{X} \mid x(k))}_{\geq p} p(x(k) \mid x(0)) dx(k) \geq p$$

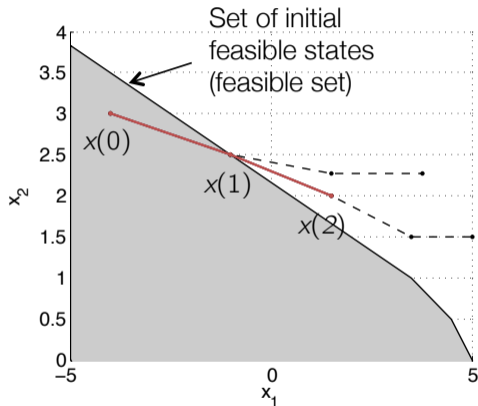
But it can be subject to considerable conservatism

Recap: Feasibility Issues in SMPC

Problem: (Potentially unbounded) stochastic disturbances can drive state initial state $x_0 = x(k)$ into infeasible region.

Several strategies to handle this problem

- Assume a maximum size of disturbance
→ use robust techniques
- Make use of recovery mechanisms
- Alternative forms of feedback

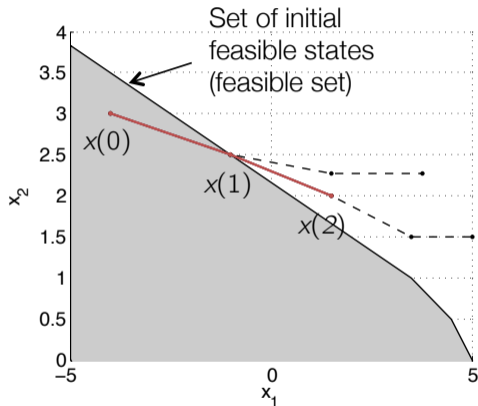


Recap: Feasibility Issues in SMPC

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Recap: Stochastic MPC with Bounded Disturbances

Setup:

$$x(k+1) = Ax(k) + Bu(k) + w(k)$$

where $w(k) \sim \mathcal{Q}^w$ i.i.d. **and** all $w(k) \in \mathcal{W}$ with \mathcal{W} a compact set.

- Stochastic disturbance $w(k)$ has **bounded support** \mathcal{W}
 - Enables use of robust techniques for recursive feasibility
-

One possible approach related to "constraint-tightening" robust MPC:

- Enforce **chance** constraints w.r.t. **all** possible previous disturbances ($i-1$ -steps robust, 1-step stochastic)
 - Enforce terminal robust invariant set (within constraints) robustly
 - For simplicity, we neglect input constraints for now, extension is straightforward
 - References: [4, 5]
-

Recap: Stochastic MPC with Bounded Disturbances

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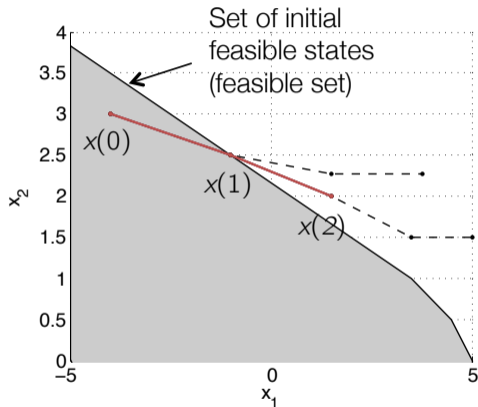
- Recursive feasibility properties follow directly from robust case
- Tractable in particular for polytopic \mathcal{W} and half-space chance constraints (or polytopic constraint as collection of half-space chance constraints)
- Asymptotic average performance analysis directly applicable (candidate solution feasible)

Feasibility Issues in SMPC (This Chapter)

Problem: (Potentially unbounded) stochastic disturbances can drive state initial state $x_0 = x(k)$ into infeasible region.

Several strategies to handle this problem

- Assume a maximum size of disturbance
→ use robust techniques
- **Make use of recovery mechanisms**
- **Alternative forms of feedback**



Contents – 7: Stochastic MPC II

Understand different strategies of dealing with feasibility issues under unbounded noise, in particular:

- Recovery mechanisms
 - implications on stability
 - implications on closed-loop constraint satisfaction guarantees
- Indirect feedback
 - implications on stability
 - implications on closed-loop constraint satisfaction guarantees
 - practical implications

Outline

1. Stochastic MPC for unbounded disturbances with "Open-Loop MPC"
2. Stochastic MPC for unbounded disturbances with "Recovery Initialization"
3. Stochastic MPC for unbounded disturbances with "Indirect Feedback"

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Recall: Chance Constrained Optimal Control

$$\begin{aligned} J^*(x) = \min_{\{u_k\}} \quad & \mathbb{E} \left(\sum_{i=0}^{\bar{N}-1} \|x(k)\|_Q^2 + \|u(k)\|_R^2 \right) \\ \text{s.t.} \quad & x(k+1) = Ax(k) + Bu(k) + w(k), \\ & u(k) = K(x(k) - \mathbb{E}(x(k))) + u_k, \\ & w(k) \sim \mathcal{N}(0, \Sigma_w), \text{ i.i.d.}, \\ & \Pr(h_{x,j}^\top x(k) \leq b_{x,j}) \geq p_{x,j} \quad \forall j = 1, \dots, N_{c,x}, \\ & \Pr(h_{u,j}^\top u(k) \leq b_{u,j}) \geq p_{u,j} \quad \forall j = 1, \dots, N_{c,u}, \\ & x(0) = x \end{aligned}$$

- (Approximate) solution to chance constrained stochastic optimal control problem
- Tractable in particular for Gaussian disturbances (unbounded support!)

Recall: Chance Constrained Optimal Control

$$\begin{aligned}\tilde{J}^*(x) = & \min_{\{\bar{u}_k\}} \sum_{k=0}^{\bar{N}-1} \|\mathbb{E}(x(k))\|_Q^2 + \|\bar{u}_k\|_R^2 \\ \text{s.t. } & \mathbb{E}(x(k+1)) = A\mathbb{E}(x(k)) + B\bar{u}_k, \\ & \text{var}(x(k+1)) = (A + BK)\text{var}(x(k))(A + BK)^T + \Sigma_w, \\ & h_{x,j}^T \mathbb{E}(x(k)) \leq b_{x,j} - \sqrt{h_{x,j}^T \text{var}(x(k)) h_{x,j}} \phi^{-1}(p_{x,j}) \quad \forall j = 1, \dots, N_{c,x}, \\ & h_{u,j}^T \bar{u}_k \leq b_{u,j} - \sqrt{h_{u,j}^T K \text{var}(x(k)) K^T h_{u,j}} \phi^{-1}(p_{u,j}) \quad \forall j = 1, \dots, N_{c,u}, \\ & \mathbb{E}(x(0)) = x, \quad \text{var}(x(0)) = 0\end{aligned}$$

- (Approximate) solution to chance constrained stochastic optimal control problem
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Can we do this in receding horizon? Kind of..

Some Notation

Assumption: Prediction process is identical to "real" process

$$x(k+1) = Ax(k) + Bu(k) + w(k),$$

$$x_{i+1} = Ax_i + Bu_i + w_i,$$

$$w(k) \sim \mathcal{N}(0, \Sigma_w) \text{ i.i.d.}$$

$$w_i \sim \mathcal{N}(0, \Sigma_w) \text{ i.i.d.}$$

To simplify, we introduce some notation for the prediction dynamics:

$$\bar{x}_i := \mathbb{E}(x_i)$$

$$\bar{u}_i := \mathbb{E}(u_i)$$

$$d_i := x_i - \bar{x}_i$$

$$\bar{d}_i := \mathbb{E}(d_i) = 0$$

$$\Sigma_i^x := \text{var}(d_i) = \text{var}(x_i)$$

resulting in \rightarrow

$$u_i = Kd_i + \bar{u}_i$$

$$\bar{x}_{i+1} = A\bar{x}_i + B\bar{u}_i$$

$$\Sigma_{i+1}^x = (A + BK)\Sigma_i^x(A + BK)^T + \Sigma_w$$

where the expectations are understood conditioned on the initial state of the prediction $i = 0$

Intuition: "Open-Loop" Stochastic MPC

$$\begin{aligned}
 \tilde{J}^*(x(k)) = \min_{\{\bar{u}_i\}} \quad & \|\bar{x}_N\|_P^2 + \sum_{i=0}^{N-1} \|\bar{x}_i\|_Q^2 + \|\bar{u}_i\|_R^2 \\
 \text{s.t.} \quad & \bar{x}_{i+1} = A\bar{x}_i + B\bar{u}_i, \\
 & \Sigma_{i+1}^x = (A + BK)\Sigma_i^x(A + BK)^T + \Sigma_w, \\
 & h_{x,j}^T \bar{x}_i \leq b_{x,j} - \sqrt{h_{x,j}^T \Sigma_i^x h_{x,j}} \phi^{-1}(p_{x,j}) \quad \forall j = 1, \dots, N_{c,x}, \\
 & h_{u,j}^T \bar{u}_i \leq b_{u,j} - \sqrt{h_{u,j}^T K \Sigma_i^x K^T h_{u,j}} \phi^{-1}(p_{u,j}) \quad \forall j = 1, \dots, N_{c,u}, \\
 & \bar{x}_N \in \bar{\mathcal{X}}_f, \\
 & \bar{x}_0 = \bar{x}_{1|k-1}, \quad \Sigma_0^x = \Sigma_{1|k-1}^x
 \end{aligned}$$

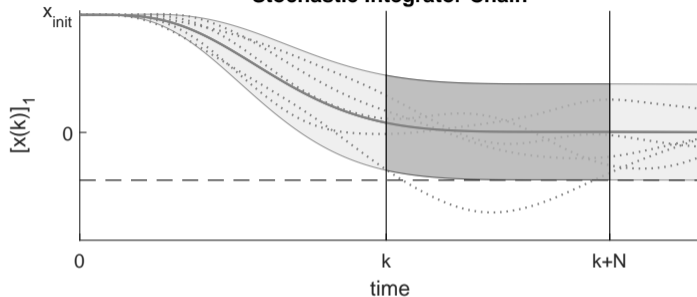
- Apply $u(k) = K(x(k) - \bar{x}_0^*) + \bar{u}_0^*$
- Initialize at predicted distribution

→ Nominal MPC: recursively feasible¹ and satisfies closed-loop chance constraints

¹Requires typical selection of terminal set for nominal MPC (within maximum tightened constraints)

Illustration: "Open-Loop" Stochastic MPC

Stochastic Integrator Chain



solid line: mean trajectory
dashed lines: trajectory samples
shaded area: confidence region

- Original optimization problem (Lecture 6) solves for entire task horizon \bar{N}
 - "Open-Loop" stochastic MPC repeatedly solves for shortened horizon N
 - Equivalent (given right terminal cost & set, sufficiently long horizon)
- **This "MPC" reduces to iterative trajectory planning method (no feedback from $x(k)$)**

Outline

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Outline

2. Stochastic MPC for unbounded disturbances with "Recovery Initialization"

- Recovery Initialization

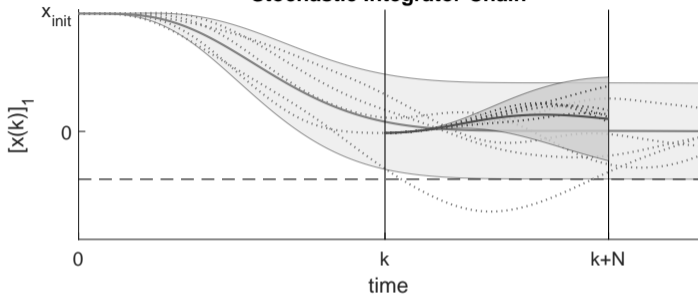
- Implications for Stability

- Implications for Chance Constraint Satisfaction

- Closed-loop Chance Constraint Satisfaction with Probabilistic Reachable Sets

Idea: Introduce Feedback whenever Feasible

Stochastic Integrator Chain



solid line: mean trajectory
dotted lines: trajectory samples
shaded area: confidence region

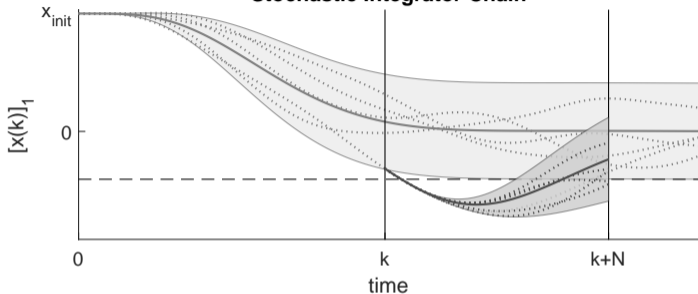
Recovery Initialization

Case 1: $\bar{x}_0 = x(k)$, $\Sigma_0^x = 0$ (Whenever feasible)

Case 2: $\bar{x}_0 = \bar{x}_{1|k-1}$, $\Sigma_0^x = \Sigma_{1|k-1}^x$ (Otherwise, guaranteed feasible)

Idea: Introduce Feedback whenever Feasible

Stochastic Integrator Chain



solid line: mean trajectory
dotted lines: trajectory samples
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Recovery Initialization

Case 1: $\bar{x}_0 = x(k)$, $\Sigma_0^x = 0$ (Whenever feasible)

Case 2: $\bar{x}_0 = \bar{x}_{1|k-1}$, $\Sigma_0^x = \Sigma_{1|k-1}^x$ (Otherwise, guaranteed feasible)

Stochastic MPC with Recovery Initialization

Aims to combine "best of both worlds":

- Feasibility guarantees from "open-loop" stochastic MPC
- Use feedback whenever possible

But theoretical analysis proves challenging!

Outlook

- What are the implications for stability/performance?
 - Candidate solution does not remain feasible (from measured state)
 - Cost from measured state may increase (over "shifted" solution) [1]
- What are implications for closed-loop constraint satisfaction?
 - No direct guarantees on closed-loop (as opposed to bounded disturbance case) [3]
 - Can also be conservative! (same for bounded disturbance case) [5]

Outline

2. Stochastic MPC for unbounded disturbances with "Recovery Initialization"

Recovery Initialization

Implications for Stability

Implications for Chance Constraint Satisfaction

Closed-loop Chance Constraint Satisfaction with Probabilistic Reachable Sets

Stability under Recovery Initialization

Recovery Initialization

Case 1 (C_1): $\bar{x}_0 = x(k), \quad \Sigma_0^x = 0$ (Whenever feasible)

Case 2 (C_2): $\bar{x}_0 = \bar{x}_{1|k-1}, \quad \Sigma_0^x = \Sigma_{1|k-1}^x$ (Otherwise, guaranteed feasible)

For asymptotic average performance bound, we want

$$\mathbb{E}(J^*(\tilde{x}(k+1) | \tilde{x}(k)) - J^*(\tilde{x}(k))) \leq -l(x(k), \pi(\tilde{x}(k))) + C$$

(Note that this is not a state-feedback controller: extended state $\tilde{x}(k) = (x(k), \bar{x}_{1|k-1}, \Sigma_{1|k-1}^x)$)

We need to show this considering both possible cases (C_1 and C_2)!

$$\mathbb{E}(J^*(\tilde{x}(k+1) | \tilde{x}(k))) = \mathbb{E}(J^*(\tilde{x}(k+1) | \tilde{x}(k), C_1) \Pr(C_1 | \tilde{x}(k))) + \mathbb{E}(J^*(\tilde{x}(k+1) | \tilde{x}(k), C_2) \Pr(C_2 | \tilde{x}(k)))$$

Case 2: Recovery Initialization

Given Case 2, we have

$$\bar{x}_0 = \bar{x}_{1|k-1}, \Sigma_0^x = \Sigma_{1|k-1}^x$$

This implies $x_0 \stackrel{d}{=} x_1^*$, in fact the candidate sequence

$$\bar{U} = \{\bar{u}_0, \dots, \bar{u}_{N-1}\} = \{\bar{u}_1^*, \dots, \bar{u}_{N-1}^*, K\bar{x}_N^*\}$$

remains feasible (shifting the previous solution works), resulting in

$$\{x_0, \dots, x_N\} \stackrel{d}{=} \{x_1^*, \dots, x_{N-1}^*, (A + BK)x_N^* + w\} \quad (\text{equal in distribution})$$

Choosing $P = (A + BK)^T P (A + BK) + Q + K^T R K$ and following the previous lecture we get

$$\mathbb{E}(J^*(\tilde{x}(k+1) | \tilde{x}(k), C_2) - J^*(\tilde{x}(k))) \leq -l(x(k), \pi(\tilde{x}(k))) + C$$

Case 1: Feedback Initialization

Given Case 1, we have

$$\bar{x}_0 = x(k), \Sigma_0^x = 0$$

and we know a feasible solution exists from this initial condition.

Problem: This does not mean that the candidate sequence is feasible!
→ shifting the previous solution **does not** work

In fact, it may happen that

$$J^*(\tilde{x}(k+1)) > \bar{J}(\tilde{x}(k+1))$$

(feasible (& optimal) solution with higher cost $J^*(\tilde{x}(k+1))$ exists)

Possible remedy [1,2]: Use feedback (C_1) only if $J^*(\tilde{x}(k+1)) \leq \bar{J}(\tilde{x}(k+1))$

$$\mathbb{E}(J^*(\tilde{x}(k+1)) \mid \tilde{x}(k), C_1) - J^*(\tilde{x}(k)) \leq -l(x(k), \pi(\tilde{x}(k))) + C$$

Stability-Adapted Recovery Scheme [1,2]

Adapted Recovery Initialization

Case 1 (C_1): $\bar{x}_0 = x(k)$, $\Sigma_0^x = 0$ (If feasible & reduces cost)

Case 2 (C_2): $\bar{x}_0 = \bar{x}_{1|k-1}$, $\Sigma_0^x = \Sigma_{1|k-1}^x$ (Otherwise, guaranteed feasible & cost decrease)

With this we have

$$\mathbb{E}(J^*(\tilde{x}(k+1) | \tilde{x}(k)) - J^*(\tilde{x}(k))) \leq -l(x(k), \pi(\tilde{x}(k))) + C$$

since

$$\begin{aligned} \mathbb{E}(J^*(\tilde{x}(k+1) | \tilde{x}(k)) - J^*(\tilde{x}(k))) &= \underbrace{\mathbb{E}(J^*(\tilde{x}(k+1) | \tilde{x}(k), C_1) - J^*(\tilde{x}(k))) \Pr(C_1 | \tilde{x}(k))}_{\leq -l(x(k), \pi(\tilde{x}(k))) + C} \\ &\quad + \underbrace{\mathbb{E}(J^*(\tilde{x}(k+1) | \tilde{x}(k), C_2) - J^*(\tilde{x}(k))) \Pr(C_2 | \tilde{x}(k))}_{\leq -l(x(k), \pi(\tilde{x}(k))) + C} \end{aligned}$$

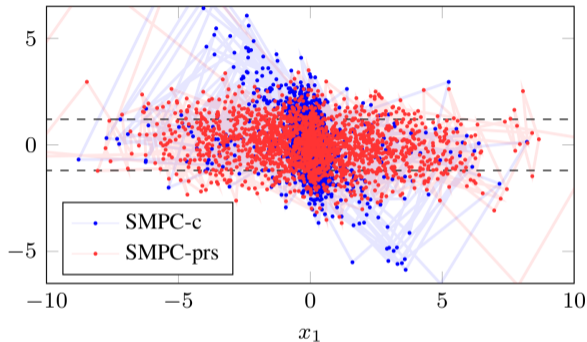
Effect of Reduced Feedback [3]

Scenario

$$x(k+1) = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} x(k) + \begin{bmatrix} 0.5 \\ 1 \end{bmatrix} u(k) + w(k)$$

Rare (unmodeled) large disturbances

- Red: C_1 whenever feasible
- Blue: C_1 feasible & cost decreases



When driven far away from the origin blue keeps candidate solution
→ no feedback on (constrained) optimization problem
→ applies mainly tube controller, ignoring constraints

Outline

2. Stochastic MPC for unbounded disturbances with "Recovery Initialization"

Recovery Initialization

Implications for Stability

Implications for Chance Constraint Satisfaction

Closed-loop Chance Constraint Satisfaction with Probabilistic Reachable Sets

Constraint Satisfaction from Recursive Feasibility

Closed-loop chance constraints

$$(*) \Pr(x(k) \in \mathcal{X} \mid x(0)) \geq p, \forall k \geq 0$$

If feasible: MPC formulation enforces (since $x_1 \stackrel{d}{=} x(k+1)$ given $x(k)$)

$$(**) \Pr(x(k+1) \in \mathcal{X} \mid x(k)) \geq p, \forall k \geq 0$$

Given feasibility at all $k \geq 0$ we have that $(**) \Rightarrow (*)$ since

$$\Pr(x(k) \in \mathcal{X} \mid x(0)) = \int \underbrace{\Pr(x(k) \in \mathcal{X} \mid x(k-1))}_{\geq p} p(x(k-1) \mid x(0)) dx(k-1) \geq p$$

Constraint Satisfaction from Recursive Feasibility

Closed-loop chance constraints

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If feasible: MPC formulation enforces (since $x_1 \stackrel{d}{=} x(k+1)$ given $x(k)$)

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Given feasibility at all $k \geq 0$ we have that $(**) \Rightarrow (*)$ since

$$\Pr(x(k) \in \mathcal{X} \mid x(0)) = \int \underbrace{\Pr(x(k) \in \mathcal{X} \mid x(k-1))}_{\geq p} p(x(k-1) \mid x(0)) dx(k-1) \geq p$$

Now we cannot assume feasibility in each time step (cannot assume $(**)$ holds):

$$\Pr(x(k+1) \in \mathcal{X} \mid x(k)) = \underbrace{\Pr(x(k+1) \in \mathcal{X} \mid x(k), C_1)}_{\geq p} \Pr(C_1) + \underbrace{\Pr(x(k+1) \in \mathcal{X} \mid x(k), C_2)}_{?? \text{ can be } \ll p} \Pr(C_2)$$

Counterexample: Closed-loop Constraint Violation I

Double integrator system:

$$x(k+1) = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} x(k) + \underbrace{\begin{bmatrix} 0.5 \\ 1 \end{bmatrix}}_B u(k) + w(k)$$

with i.i.d. Gaussian disturbance

$$w(k) \sim \mathcal{N}(0, BB^T)$$

Velocity chance constraint

$$\Pr([x(k)]_2 \geq -1.2) \geq 0.8$$

Controlled with MPC

$$\min_{\{\bar{u}_i\}} \|\bar{x}_N\|_P^2 + \sum_{i=0}^{N-1} \|\bar{x}_i\|_Q^2 + \|\bar{u}_i\|_R^2$$

$$\text{s.t. } \bar{x}_{i+1} = A\bar{x}_i + B\bar{u}_i,$$

$$\Sigma_{i+1}^x = (A + BK)\Sigma_i^x(A + BK)^T + \Sigma_w,$$

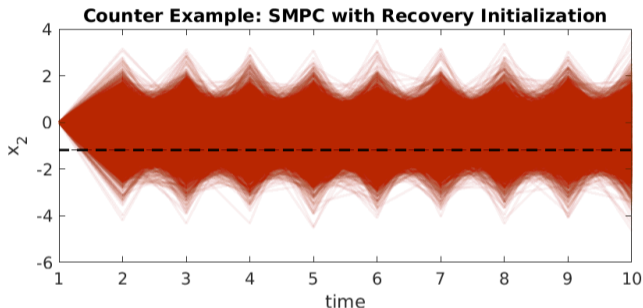
$$[\bar{x}_i]_2 \geq -1.2 + \sqrt{[\Sigma_i^x]_{2,2}} \phi^{-1}(0.8),$$

$$\bar{x}_N \in \mathcal{X}_f,$$

$$\bar{x}_0 = \begin{cases} x(k), & \Sigma_0^x = 0 \text{ feasible \& cost decrease} \\ \bar{x}_0 = \bar{x}_{1|k-1}, & \Sigma_0^x = \Sigma_{1|k-1}^x, \text{ otherwise} \end{cases}$$

$$\text{Applied input: } u(k) = K(x(k) - \bar{x}_0^*) + u_0^*$$

Counterexample: Closed-loop Constraint Violation II



- Optimal to minimize $[x(k)]_2$ as much as possible
- Given $[x(k)]_2 \ll -1.2$: MPC may be infeasible and tube controller is applied (recovery)
- Given $[x(k)]_2 \gg -1.2$: MPC aggressively drives system back into constraint (cost optimal)

Closed-loop system more 'aggressive' than open-loop plan, constraint violation: 24.1% > 20%

Outline

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Recovery Initialization

Implications for Stability

Implications for Chance Constraint Satisfaction

Closed-loop Chance Constraint Satisfaction with Probabilistic Reachable Sets

Stochastic MPC with Probabilistic Reachable Sets

Problem: The distribution of closed-loop $x(k)$ is very hard to establish.

Define $x(k) = \bar{x}_{0|k} + d(k)$. Then

- $d(k) = 0$ whenever MPC feasible, otherwise $d(k+1) = (A + BK)d(k) + w(k)$
- depends on MPC optimization

→ hard to derive constraint tightening that applies in closed-loop

Idea: We can analyze the distribution of $d_{k+1} = (A + BK)d_k + w_k$. Can we relate $d(k)$ to d_k , for $d(0) = d_0 = 0$?

General closed-loop constraint satisfaction is an open problem:

We will discuss a result for the special case of unimodal $w(k)$ and symmetric constraint tightening.

Convex Unimodal Random Variable

A random variable w is convex unimodal if for every direction $\hat{w} \in \mathbb{R}^n$ and every convex set \mathcal{F} symmetric around the origin the probability $\Pr(w + c\hat{w} \in \mathcal{F})$ is non-increasing in $c \geq 0$.

Symmetric Tightening with Unimodal Disturbances

[3, Theorem 3]

Let $w(k)$ and w_k be i.i.d. convex unimodal random variables and \mathcal{F} a convex symmetric set around the origin. Let $d_{k+1} = (A + BK)d_k + w_k$, $d_0 = 0$ be the (linearly) predicted deviation, and $d(k) = x(k) - \bar{x}_{0|k}$ the actual deviation under stochastic MPC with recovery initialization. Then

$$\Pr(d(k) \in \mathcal{F}) \geq \Pr(d_k \in \mathcal{F}) \quad \forall k \geq 0$$

Proof idea:

$$\begin{aligned} & \Pr((A + BK)d(k) + w(k) \in \mathcal{F}) \\ &= \Pr((A + BK)d(k) + w(k) \in \mathcal{F} \mid C_1) \Pr(C_1) + \Pr((A + BK)d(k) + w(k) \in \mathcal{F} \mid C_2) \Pr(C_2) \\ &= \Pr(w(k) \in \mathcal{F} \mid C_1) \Pr(C_1) + \Pr((A + BK)d_k + w(k) \in \mathcal{F} \mid C_2) \Pr(C_2) \\ &\geq \Pr((A + BK)d(k) + w(k) \in \mathcal{F} \mid C_1) \Pr(C_1) + \Pr((A + BK)d_k + w(k) \in \mathcal{F} \mid C_2) \Pr(C_2) \\ &= \Pr((A + BK)d_k + w_k \in \mathcal{F}) \end{aligned}$$

(due to unimodality of $w(k)$ and symmetry of \mathcal{F})

Probabilistic Reachable Sets

This result can be used to derive sets for constraint tightening that hold in closed-loop:

Probabilistic Reachable Set [3, 4]

A set \mathcal{F}_i^p is an i -step probabilistic reachable set (PRS) of probability level p for process $\{d(k)\}$ initialized at $d(0)$ if

$$\Pr(d(i) \in \mathcal{F}_i^p \mid d(0)) \geq p$$



Reformulation of chance constraints with $x(k) = \bar{x}_0 + d(k)$ and \mathcal{F}_k^p PRS for $d(k)$.

$$\Pr(x(k) \in \mathcal{X}) \geq p \Leftrightarrow \bar{x}_0 \in \mathcal{X} \ominus \mathcal{F}_k^p$$

([3, Theorem 3]) PRS for $d_{k+1} = (A + BK)d_k + w_k$ is also PRS for $d(k) \rightarrow$ tractable computation

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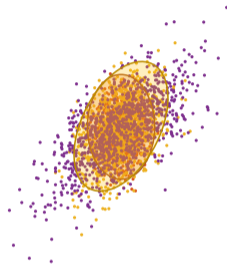
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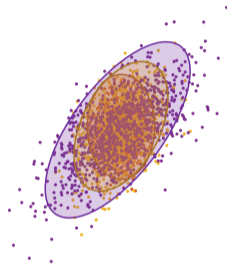
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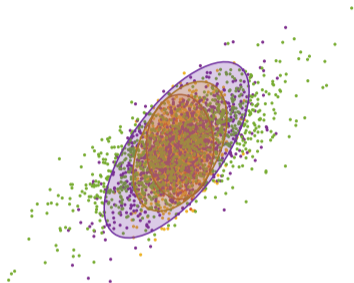
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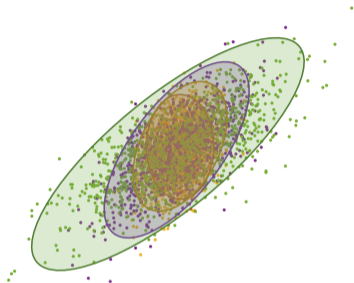
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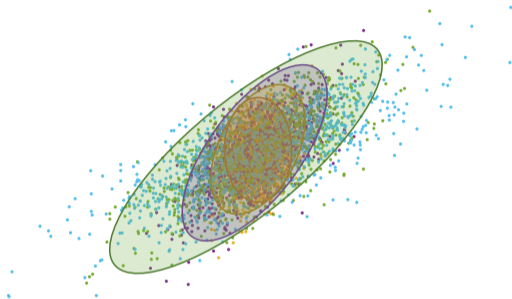
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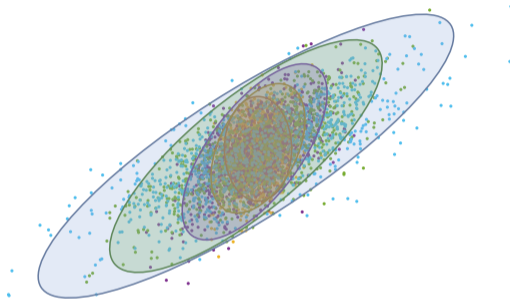
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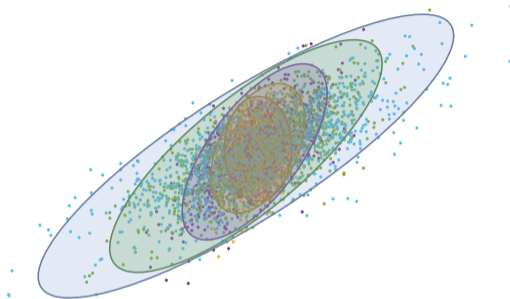
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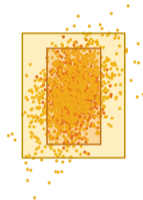
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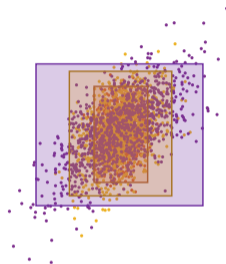
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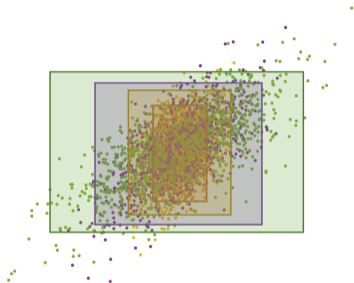
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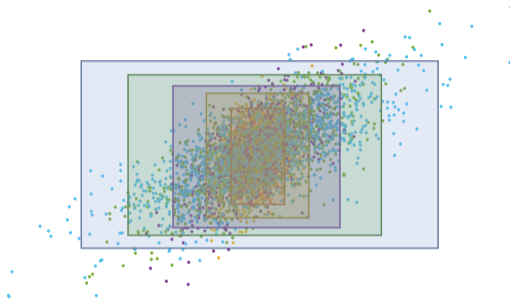
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SMPC with PRS & Recovery Initialization

$$\begin{aligned} \min_{\{\bar{u}_i\}} \quad & \|\bar{x}_N\|_P^2 + \sum_{i=0}^{N-1} \|\bar{x}_i\|_Q^2 + \|\bar{u}_i\|_R^2 \\ \text{s.t.} \quad & \bar{x}_{i+1} = A\bar{x}_i + B\bar{u}_i, \\ & \bar{x}_i \in \mathcal{X} \ominus \mathcal{F}_{i+k}^p, \\ & \bar{x}_N \in \mathcal{Z}_f \subset \mathcal{X} \ominus \mathcal{F}_\infty^p, \\ & \bar{x}_0 = \begin{cases} x(k), & \text{if feasible} \\ \bar{x}_{1|k-1}, & \text{otherwise} \end{cases} \end{aligned}$$

Closed-loop constraint satisfaction guarantee if $w(k)$ unimodal and \mathcal{F}_i^p symmetric

Summary

- Result for tightening with convex symmetric probabilistic reachable sets
- Alternatives exists (e.g. soft constraints) but analysis similarly lacking

Outline

1. Stochastic MPC for unbounded disturbances with "Open-Loop MPC"
2. Stochastic MPC for unbounded disturbances with "Recovery Initialization"
3. Stochastic MPC for unbounded disturbances with "Indirect Feedback"

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3. Stochastic MPC for unbounded disturbances with "Indirect Feedback"
Constraint Satisfaction in Prediction

Chance Constraints in Closed-Loop vs. in Prediction

We want to show:

$$(*) \Pr(x(k) \in \mathcal{X} \mid x(0)) \quad (\text{conditioned on initial state } x(0))$$

But typically enforce

$$(**) \Pr(x_i \in \mathcal{X} \mid x_0 = x(k)) \quad (\text{conditioned on measured state } x(k))$$

Goal was to guarantee that $(**)$ always holds, and then $(**) \Rightarrow (*)$

- Constraint tightening SMPC for bounded disturbances (last lecture)

In general, enforcing $(**)$ for all k is much stricter than $(*)$

- Leads to feasibility issues
- Can be conservative, if $(*)$ is the required condition

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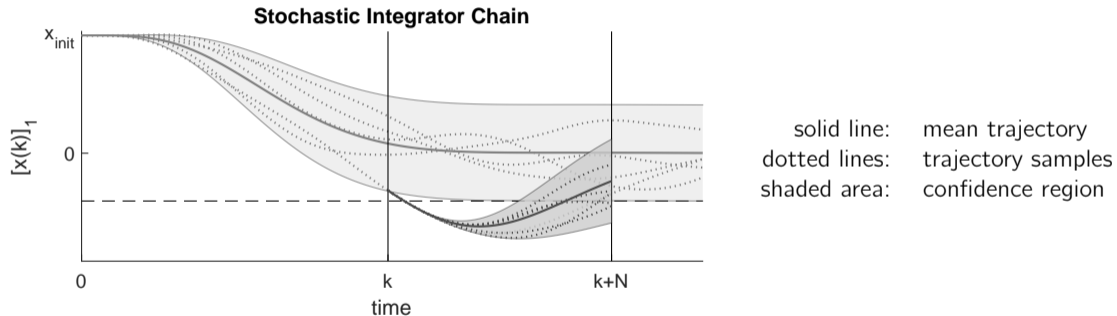
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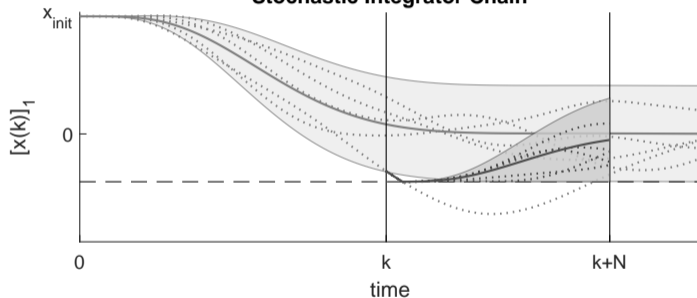
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Chance Constraints in Closed-Loop vs. in Prediction



Chance Constraints in Closed-Loop vs. in Prediction

Stochastic Integrator Chain



solid line: mean trajectory
dotted lines: trajectory samples
shaded area: confidence region

Chance Constraints in Closed-Loop vs. in Prediction

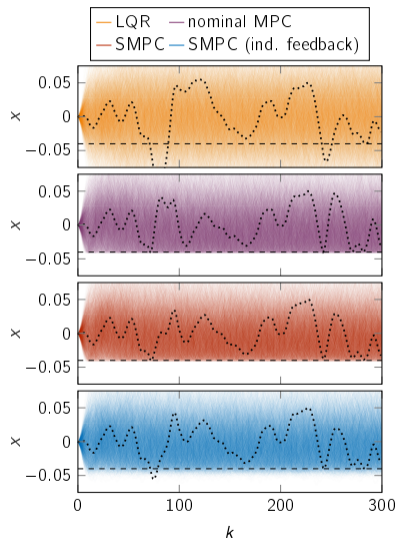
Third order integrator chain:

$$x(k+1) = \begin{bmatrix} 1 & 0.1 & 0.1^2/2 \\ 0 & 1 & 0.1 \\ 0 & 0 & 1 \end{bmatrix} x(k) + \begin{bmatrix} 0.1^3/6 \\ 0.1^2/2 \\ 0.1 \end{bmatrix} (u(k) + w(k))$$

with i.i.d. $w(k) \sim \mathcal{N}(0, 1)$ and chance constraint

$$\Pr([x]_1 \geq -\sqrt{[\Sigma_\infty]_{1,1}}) \geq 0.84$$

- Unconstrained (LQR) solution satisfies constraint
- Main effect of disturbance on $[x]_1$ is delayed (needs to propagate through system)
- SMPC results in virtually no "constraint softening"
- Side effect: Aggressive control inputs to ensure feasibility



Indirect feedback SMPC: Idea

Introduce **nominal state** $z(k)$ and let it evolve according to **nominal input** $v(k) = v_0^*$

$$z(k+1) = Az(k) + Bv(k)$$

$$x(k+1) = Ax(k) + Bu(k) + w(k) = Ax(k) + B(K(x(k) - z(k)) + v(k)) + w(k)$$

$$e(k+1) = x(k+1) - z(k+1) = (A + BK)e(k) + w(k)$$

→ no **direct feedback** from measurement $x(k)$ on $z(k)$

→ error state $e(k) = x(k) - z(k)$ evolves linearly and independent of $v(k)$

Straightforward to formulate constraints on closed-loop (similar to 'open-loop' stochastic MPC)

$$z(k) \in \mathcal{X} \ominus \mathcal{F}_k^p \Rightarrow \Pr(x(k) \in \mathcal{X}) \geq p$$

As opposed to "open-loop" stochastic MPC, we nevertheless introduce feedback by optimizing over cost given measured state $x(k)$, i.e. $x_0 = x(k)$ (**indirect feedback**)

$$\min \mathbb{E} \left(\left\| x_N \right\|_P^2 + \sum_{i=0}^{N-1} \left\| x_i \right\|_Q^2 + \left\| u_i \right\|_R^2 \middle| x_0 = x(k) \right)$$

Indirect Feedback SMPC: Resulting Formulation

$$\begin{aligned}\tilde{J}^*(x(k)) = \min_{\{v_i\}} \quad & \|\bar{x}_N\|_P^2 + \sum_{i=0}^{N-1} \|\bar{x}_i\|_Q^2 + \|\bar{u}_i\|_R^2 \\ \text{s.t.} \quad & \bar{x}_{i+1} = A\bar{x}_i + B\bar{u}_i, \\ & z_{i+1} = Az_i + Bv_i, \\ & \bar{e}_i = \bar{x}_i - z_i, \\ & \bar{u}_i = K\bar{e}_i + v_i, \\ & z_i \in \mathcal{X} \ominus \mathcal{F}_{k+i}^p, \\ & z_N \in \mathcal{Z}_f \subset \mathcal{X} \ominus \mathcal{F}_\infty^p, \\ & \bar{x}_0 = x(k), \quad z_0 = z(k) = z_{1|k-1}\end{aligned}$$

- Applied input: $u(k) = K(x(k) - z(k)) + v_0^*$
- \mathcal{Z}_f terminal (nominal) invariant set
- \mathcal{F}_{k+i}^p probabilistic reachable sets for $e(k+i)$

Indirect Feedback SMPC: Main Properties

Candidate sequence remains feasible (details in recitation)

$$\begin{aligned}\bar{V} &= \{v_1^*, \dots, v_{N-1}^*, \pi_f(z_N^*)\} \\ \rightarrow \bar{Z} &= \{z_1^*, \dots, z_N^*, Az_N^* + B\pi_f(z_N^*)\}\end{aligned}$$

From this follows recursive feasibility and

- Closed-loop chance constraint satisfaction, since

$$x(k) = z(k) + e(k), \quad z(k) \in \mathcal{X} \ominus \mathcal{F}_k^p \text{ and } \Pr(e(k) \in \mathcal{F}_k^p) \geq p$$

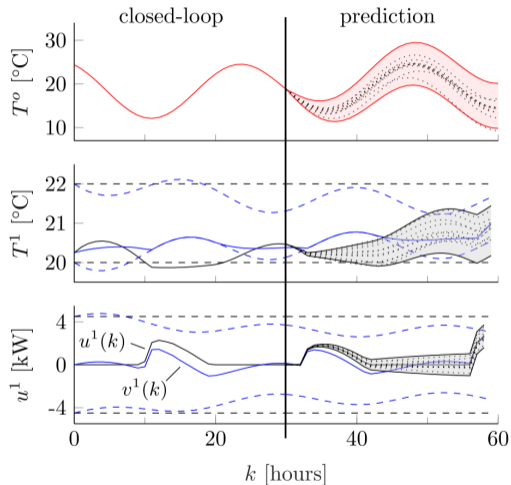
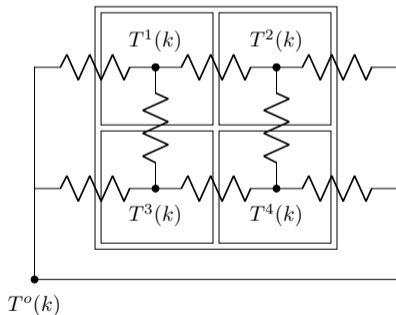
- Asymptotic average performance bound

$$\lim_{\bar{N} \rightarrow \infty} \frac{1}{\bar{N}} \sum_{k=0}^{\bar{N}-1} \mathbb{E}(l(x(k), u(k))) \leq \text{tr}(P\Sigma_w)$$

when choosing $\pi_f(x) = Kx$ and $P = (A + BK)^\top P(A + BK) + Q + K^\top RK$

Indirect Feedback SMPC: Building Control

- Modeled as resistance network (linear)
- Heating of 4 different Rooms
- Comfort constraint $\Pr(T_j \in [20, 22]) \geq 0.9$
- Sparsity inducing input cost $\|u\|_1$ (1-norm)



Indirect Feedback SMPC: Practical Effects

Standard double integrator system:

$$x_{i+1} = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} x_i + \underbrace{\begin{bmatrix} 0.5 \\ 1 \end{bmatrix}}_B u_i + w_i, \quad w_i \sim \mathcal{N}(0, BB^T)$$

Velocity chance constraint $\Pr([x(k)]_2 \leq 3) \geq 0.9$

But now: **Model mismatch** $B_{\text{sim}} = \frac{1}{5}B$

| Controller | Closed-Loop Cost | Constraint Violation |
|--------------------------|------------------|----------------------|
| Nominal MPC | 130 | 21.7% |
| (direct feedback) SMPC | 145 | 11.0% |
| (indirect feedback) SMPC | 91 | 19.8% |

Summary

- Stochastic MPC remains active research topic, in particular
 - unbounded disturbance distributions (this lecture)
 - nonlinear stochastic MPC (not discussed)
- Common formulation (satisfying chance constraints in prediction) may need critical re-evaluation
 - Indirect feedback as possible alternative
 - Connection to stochastic reference governors

References and further reading

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